



## **HPC Challenges and New Computing Frontiers**

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Workshop on HPC challenges for new extreme scale applications Paris, March 6<sup>th</sup>, 2023

### Preamble

I survey some results obtained for sparse linear algebra for iterative methods and for machine learning methods.

I also discuss on the potential evolution we would face to be able to mix computational science, data science and machine learning on future faster supercomputers, based on some published examples.

#### <u>Workshop</u>

End-users and scientists have to face a lot of challenge associated to these evolutions and the increasing size of the data. The convergence of Data science (big Data) and the computational science to develop new applications generates important challenges.

## Outline

- Introduction
- New levels of programming (Graphs of Tasks, Network on chip)
- New methods and algorithms (Unite&Conquer, Stochastic Matrix,..)
- HPC and Machine Learning (GCN, Transformer,..)
- Generators of Data Sets and matrices for brain-scale applicatons
- What post-exascale plateforms and programin paradigms

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Computational science : IEEE 64 bit arihmetic,

exascale supercomputers, C++, CSR-ELLPACK

**BigData** : MapReduce, exascale 16-32 bits supercomputers, COO, Python, Data Center, (HPC on) Cloud,

Machine Learning : Diferent arithmetics (16, 32, not always IEEE), Tensors, Python

• Mainly developped by GAFAM and BATX, Nvidia and others

#### Exascale supercomputers are now availables

- Frontiers,...
- Sunway,...
- Cerebras,...

# Nevertheless, the target applications used different arithmetics and programming paradigms, and **only a few aplications reach the exascale (HPCG : 16 Pflops, Fugaku)**

Machine learning and AI applications are now requiring exascale machines, which were not first designed for them. New machines and processors (and the next generation of. post-exascale machines) are (would) be targeting "mainly" these applications.

#### The requiered arithmetic, data structures, linear algebra are often diferents.

The most expensive (time, energy) are the data migrations and communications, especially the I/O : Distributed and Parallel computing where are the data (HPC on Cloud or on DataCenter), or **generation of the data in Parallel**.

Back to

- "true" data parallelism (history : Connection Machines): Cerebras
- data flow programing (history : the Arwind MIT data flow machine): SambaNova

### We have to experiments on "new" methods and propose the generation of "brainscale" data sets (graphs-matrices) for computational science and machine learning applications.

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## 1 - Network on Chip

and irregular "local" communications

Larger number of nodes (task programming)

Network on chip :

Distributed and data parallelism

High hiearachy of execution models, which lead to several programming paradims for a given method : graph of tasks, PGAS, data parallel

#### Experiments :

- 1 MPI "task" per chip (1 openMP per chip + SVE)
- 4 MPI "tasks" per chip (1 openMP per CMG + SVE)

2021 IEEE International Conference on Cluster Computing (CLUSTER)

### Sequences of Sparse Matrix-Vector Multiplication on Fugaku's A64FX processors

Jérôme Gurhem\*, Maxence Vandromme\*, Miwako Tsuji $^{\dagger}$ , Serge G. Petiton\* $^{\ddagger}$ , Mitsuhisa Sato $^{\dagger}$ 



Sunway





Fig. 1. C-diagonal Q-perturbed sparse matrix with C=4 and Q=0.05 on the left, Q=0.5 on the middle and Q=0.9 on the right

Fig. 8. A(Ax + x) + x with OpenMP for C-diagonal Q-perturbed matrices with C = 16

## PageRank on Fugaku

thraada	C	0.0	0	.4	0.8			
uneaus	1	2	1	2	1	2		
1	32.5	35.7	9.4	15.3	6.2	10.7		
2	32.2	36.3	9.4	15.3	6.2	10.8		
4	35.8	55.4	9.7	18.2	6.6	12.2		
6	42.5	55.0	10.2	18.2	7.0	12.3		
12	51.7	64.2	11.2	19.8	8.7	13.6		
24	72.4	81.3	8.1	14.4	7.7	13.2		
48	50.3	103.9	2.7	9.8	1.3	5.2		
MPI	32.5	33.6	9.6	14.5	6.2	10.4		
			TA	TABLE VIII				

MPI AND MPI+OPENMP PERFORMANCE (IN GFLOP/S) FOR A(Ax + x) + x with C-diagonal Q-perturbed matrices on 1 and 2 Nodes for CSR storage format with different number of threads per MPI process, keeping 48 threads per node

throada	C	).0	0	.4	0.8		
unreads	1	2	1	2	1	2	
1	32.5	35.2	9.5	15.3	6.2	10.7	
2	31.9	36.8	9.5	15.6	6.2	10.9	
4	35.4	53.8	9.9	18.5	6.5	12.2	
6	41.5	55.3	10.3	18.6	6.7	12.4	
12	51.3	59.3	11.4	20.3	8.7	13.0	
24	71.8	83.3	8.2	14.7	7.2	10.5	
48	50.3	106.1	2.9	10.3	1.3	5.6	
MPI	32.6	33.1	9.6	14.8	6.3	10.4	
			TA	BLE IX			

MPI AND MPI+OPENMP PERFORMANCE (IN GFLOP/S) FOR A(Ax + x) + x with C-diagonal Q-perturbed matrices on 1 and 2 NODES FOR ELL STORAGE FORMAT WITH DIFFERENT NUMBER OF THREADS PER MPI PROCESS, KEEPING 48 THREADS PER NODE March 6, 2023 HPC challenges

### Sequences of Sparse Matrix-Vector Multiplication on Fugaku's A64FX processors

Jérôme Gurhem\*, Maxence Vandromme\*, Miwako Tsuji<sup>†</sup>, Serge G. Petiton\*<sup>‡</sup>, Mitsuhisa Sato<sup>†</sup>

#### N = 4 000 000, NNZ = 50, or 100 per node

1 MPI\_OpenMP/node better if the Matrices are not "too" iregular

Otherwise, 4 MPI-OpenMP per node are more efficient.

nodos	$\operatorname{CSR}$		E	$\operatorname{ELL}$		00	nodoa		$\mathbf{CSR}$	
nodes	node	CMG	node	CMG	node	CMG	nodes	node	CMG	noc
1	1.89	0.91	1.30	0.91	5.19	2.17	1	5.90	1.28	2.5
<b>2</b>	2.17	0.76	1.41	0.79	4.24	2.01	<b>2</b>	3.92	1.15	2.4
4	1.98	0.69	1.33	0.71	3.28	1.84	<b>4</b>	3.14	0.90	1.8
8	1.58	0.54	1.02	0.55	2.57	1.47	8	2.75	0.78	1.7
16	1.39	0.47	0.98	0.48	2.24	1.28	16	2.78	0.77	1.6
32	1.39	0.46	0.88	0.54	2.25	1.33	32	2.77	0.77	1.8
64	1.40	0.46	0.93	0.47	2.24	1.32	64	2.38	0.67	2.2
128	1.20	0.43	1.22	0.40	1.89	1.10	128	1.98	0.56	2.0
256	1.00	0.36	1.02	0.35	1.56	0.95	256	1.99	0.56	2.0
512	1.00	0.35	1.00	0.35	1.13	0.84	512	1.96	0.56	1.9
1024	1.00	0.37	1.01	0.37	1.16	0.86	1024	1.98	0.59	1.9
Table	1: N	Media	n rui	ntime	for	the	Table	2: 1	Media	n r
PageF	Rank,	$\operatorname{scali}$	ng tl	PageF	lank,	$\operatorname{scali}$	ng			
nonze	ro ele	ement	s per	nonze	ro ele	ement	s pe			
numb	er of	com	oute 1	numb	er of	com	oute			
base of	of nnz	z = 50	)	base c	of nnz	$z = 10^{10}$	00			

nodoa	$\mathbf{C}$	$\operatorname{SR}$	$\mathbf{E}$	$\operatorname{LL}$	SCOO		
noues	$\operatorname{node}$	CMG	$\operatorname{node}$	CMG	node	CMG	
1	5.90	1.28	2.59	1.30	5.59	3.53	
2	3.92	1.15	2.46	1.19	4.55	3.24	
4	3.14	0.90	1.89	0.92	4.37	2.59	
8	2.75	0.78	1.70	0.79	3.77	2.27	
16	2.78	0.77	1.69	0.81	3.83	2.27	
32	2.77	0.77	1.83	0.81	3.81	2.31	
64	2.38	0.67	2.20	0.70	3.26	1.99	
128	1.98	0.56	2.00	0.58	2.68	1.63	
256	1.99	0.56	2.00	0.56	2.68	1.64	
512	1.96	0.56	1.96	0.57	2.26	1.55	
1024	1.98	0.59	1.97	0.59	2.28	1.58	
Table	9. N	India		atima	for	the	

runtime for the the number of er row with the nodes, from a

## 2 - Graph of Tasks programming

Other new level on programming : graph of task programming



HPC challenges

YML (since 2000) - yml.prism.uvsq.fr (opensource, Cecil Licence)

We have a Virtual Machine with tutorials

# <u>yml.prism.uvsq.fr</u>

Main properties :

- High level graph description language (*coordination/control language*) *LL(1) grammar*
- Independent of Middleware, hardware and libraries
- A backend for each systems or middleware (*then platforms or supercomputers/hypercomputers*) : Xtremweb(P2P), OmniRPC, Xtremweb-OmniRPC
- Expertise may be proposed by end-users
- May use existing components / thought eventualy libraries

Deployed in France, Belgium, Ireland, Japan (K, T2K-Tsukuba, FX10-AICS) China (Hohai, Najing), Tunisia, USA (Hooper-LBNL, TOTAL-Houston). Experiment on P2P or GRID platforms : Grid (Gird5000) and P2P (100 PCs in Lille, 100 PC in France, and 4 clusters in Japan, launch from a SC INRIA booth

## Graph (n dimensions) of components/tasksYML



par compute tache1(..); notify(e1); // compute tache2(..); migrate matrix(..); notify(e2);

wait(e1 and e2); Par (i :=1;n) do par compute tache3(..); notify(e3(i)); // if(l < n)then wait(e3(i+1)); compute tache4(..); notify(e4); endif; //

compute tache5(..); control robot(..); notify(e5); violate meen(...); end par end do par // wait(e3(2:n) and e4 and e5); compute tache6(..); .../... end par

# **Multi-Level Parallelism Integration:** YML-XMP

#### OpenMP <TASK 1> NODE NODE NODE GPGPU etc... NODE NODE NODE <TASK 4> <TASK 2> <TASK 3> for(i=0;i<n;i++)<TASK 6> <TASK 5> for(j=0;j<n;j++){ tmp[i][j]=0.0; #pragma xmp loop (k) on t(k) for(k=0;k<n;k++){ tmp[i][j]+=(m1[i][k]\*m2[k][j]); <TASK 7> }}} #pragma xmp reduction (+:tmp) YML provides a workflow programming environment and high level graph description language called YvetteML

N dimension graphs available

Each task is a parallel program over several nodes. XMP language can be used to descript parallel program easily!

YML/XMP/StarPu expriments on T2K in Japan, French-Japanese project FP3C

# Experiments (2) BGJ on K-Computer



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### 1 - Minimizing the number of operations Efficient Parallel PageRank Algorithm for Network Analysis

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ParSoc22, proceedings of IPDPS22

Abstract—We propose an efficient version of the PageRank algorithm for adjacency matrices, that reduces the complexity by a factor two. This method computes the  $A^T x$  operation on the transpose matrix  $A^T$  without having to explicitly normalize and transpose the matrix. We implement the method using standard row-major and column-major matrix storage formats. We perform experiments with parallel implementations in OpenMP, on synthetic data as well as on matrices extracted from large-scale graphs. The experiments are done on two different Intel processors from recent generations. The columnmajor storage format version of our method shows good scaling and outperforms the standard PageRank in a majority of cases, even when not considering the preprocessing burden in the latter.

#### Based on the optimisation of the numer of operations for stochastic matrix by a vector products

TABLE II: Median runtime (in ms) of all three applications on synth-3, for both storage formats and both variants of SpMV (o = original, n = new) on Ruche

		Spl	MV		~	A(Aa	(x + x)			Page	Rank	
threads	CSR-o	CSR-n	CSC-o	CSC-n	CSR-o	CSR-n	CSC-o	CSC-n	CSR-o	CSR-n	CSC-o	CSC-n
1	259.9	197.1	302.6	200.8	533.8	405.5	607.8	398.3	1611	1253	1788	1186
2	131.1	112.6	151.3	100.5	264.3	225.3	310.8	199.7	807.7	673.6	905.8	594.7
4	67.2	56.3	77.9	50.6	135.2	113.1	158.0	100.1	413.7	339.2	461.4	298.8
8	35.1	28.5	41.3	25.5	70.1	57.4	81.2	50.6	215.7	172.4	266.8	151.3
12	24.2	19.4	28.5	17.2	48.4	39.1	56.0	34.1	149.3	117.2	172.5	102.3
16	18.7	14.8	22.3	13.0	37.6	29.7	44.0	25.8	116.1	90.7	133.3	77.7
20	15.5	12.1	19.1	10.5	31.1	24.6	37.9	20.9	96.2	74.6	115.4	63.2
30	14.7	10.2	17.1	7.9	29.7	21.2	34.4	15.5	92.1	65.8	102.7	46.8
40	17.2	9.5	19.5	6.6	34.8	22.1	39.0	13.0	108.1	67.7	119.3	39.8

TABLE III: Median runtime (in ms) of all three applications on synth-3, for both storage formats and both variants of SpMV (o = original, n = new) on Ice Lake

		Spl	MV			A(Ax)	(x + x)			Page	Rank	
threads	CSR-o	CSR-n	CSC-o	CSC-n	CSR-o	CSR-n	CSC-o	CSC-n	CSR-o	CSR-n	CSC-o	CSC-n
1	138.2	109.9	142.9	119.7	278.3	220.9	284.4	239.7	857.4	674.5	900.0	718.7
2	75.2	57.1	79.5	60.6	152.3	114.6	158.1	121.7	484.6	339.5	483.1	366.2
4	38.2	28.3	40.4	30.4	80.3	57.8	81.6	61.4	254.2	168.3	260.6	186.6
8	19.7	14.3	23.2	15.2	40.1	30.1	41.9	31.1	126.1	90.0	133.4	92.1
12	13.4	9.9	18.0	10.2	28.1	20.9	32.7	20.8	84.9	60.8	106.9	62.7
20	11.4	6.7	11.2	6.2	18.6	13.8	21.3	12.6	57.2	41.7	66.1	38.9
28	8.5	5.7	9.1	4.7	15.8	11.5	18.7	9.4	47.2	34.2	53.9	28.9
38	7.5	5.6	9.2	3.6	14.6	10.9	18.4	7.2	42.0	33.2	47.3	22.9
57	7.4	6.1	9.7	2.7	15.2	12.6	19.7	5.5	51.9	37.8	53.5	17.7
76	10.8	6.7	16.1	3.2	21.9	14.3	25.1	6.6	41.8	46.7	65.7	22.2

Mai

### 2 - Unite and Conquer (asynchronous computation to minimize the number of iterations)



March 6, 2023

SIAM Journal on Scientific Computing  $\rightarrow$  Vol. 27, Iss. 1 (2005)  $\rightarrow$  10.1137/S1064827500366082 Multiple Explicitly Restarted Arnoldi Method for Solving Large Eigenproblems Nahid Emad, Serge Petiton, and Guy Edjlali

## Asynchronous Iterative Restarted Methods

Collaboration with He Haiwu and Guy Bergère (U. Lille 1, CNRS) and Ye Zhang (Hohai Univ. Nanjing), Salim Nahi (Maison de la simulation), and Pierre-Yves Aquilenti (TOTAL)



- Haiwu He, Guy Bergère, and Serge Petiton, Computational Math. Appl., 2006
- Ye Zhang, Guy Bergère, and Serge Petiton, LNCS, Springer Verlag, 2008
- .../...

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#### Parallel Jaccard and Related Graph Clustering Techniques

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#### **6 GRAPH CLUSTERING**

In graph clustering a vertex set V is often partitioned into p disjoint sets  $S_k$ , such that  $V = S_1 \cup S_2 \cup S_p$  and  $S_i \cap S_j = \{\emptyset\}$  for  $i \neq j$ [16, 21]. Notice that instead of the original graph G = (V, E) we can use the modified graph  $G^{(*)} = (V^{(*)}, E^{(*)})$ , with vertex  $v_i^{(*)}$  and edge  $w_{ij}^{(*)}$  weights computed based on PageRank and Jaccard or related schemes discussed in earlier sections.

#### 6.1 Jaccard Spectral Clustering

Notice that we can define the Laplacian as

$$L^{(*)} = D^{(*)} - A^{(*)}$$
(32)

where  $D^{(*)} = \text{diag}(A^{(*)}\mathbf{e})$  is the diagonal matrix.

Then, we would minimize the normalized balanced cut

$$\tilde{\eta}(S_1, \dots, S_p) = \min_{\substack{S_1, \dots, S_p \\ U^T D^{(*)} U = I}} \sum_{k=1}^p \frac{\operatorname{vol}(\partial(S_k))}{\operatorname{vol}(S_k)}$$

$$= \min_{\substack{U^T D^{(*)} U = I}} \operatorname{Tr}(U^T L^{(*)} U)$$
(33)

where Tr(.) is the trace of a matrix, boundary edges

$$\partial S = \{(i,j) \mid i \in S \land j \notin S\}$$
(34)

and volume

$$\operatorname{vol}(S) = \sum_{i \in S} w_{ij}^{(*)}$$
(35)

$$\operatorname{vol}(\partial S) = \sum_{(i,j)\in\partial(S)} w_{ij}^{(*)} = \sum_{(i,j)\in\partial(S)} w_{ij}^{(O)} \left(1 + \frac{w_{ij}^{(I)}}{w_{ij}^{(U)}}\right)$$

by finding its smallest eigenpairs and transforming them into assignment of nodes into clusters [22]. Notice that Jaccard weights correspond to the last term in the above formula, and are related to the sum of ratios of the intersection and union of nodes on the boundary of clusters.







Figure 2: Amazon book co-purchasing graph with Jaccard

**HPC** challenges

## Graph Convolutional Network (GCN)

#### Enhancing Graph Convolutional Networks by Topology Sampling

# $H^{(\ell+1)} = \sigma (\tilde{A} H^{(\ell)} W^{(\ell)}); \ell = 0, L^{*}$

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Normalized Laplacian (Sparse non symetrical) Matrices (computed from the adjacency matrix)

**Sequence of sparse by dense by dense products** 

Edge and/or node droping to limit the over-smoothing



Dropped Subgraph

### Transformer method general structure



BERT, GPT,..

#### Attention :



# Sequence of dense by dense rectangular matrix products

Brain scale experiment : BaGuaLu (China, 1 exascale, mixed arithmetic, Sunway – processor, with a network on chip) Proceedings of PPoPP 22

#### BAGUALU: Targeting Brain Scale Pretrained Models with over 37 Million Cores

Zixuan Ma<sup>1</sup>, Jiaao He<sup>1</sup>, Jiezhong Qiu<sup>1,4</sup>, Huanqi Cao<sup>1</sup>, Yuanwei Wang<sup>1</sup>, Zhenbo Sun<sup>1</sup>, Liyan Zheng<sup>1</sup>, Haojie Wang<sup>1</sup>, Shizhi Tang<sup>1</sup>, Tianyu Zheng<sup>3</sup>, Junyang Lin<sup>2</sup>, Guanyu Feng<sup>1</sup>, Zeqiang Huang<sup>3</sup>, Jie Gao<sup>3</sup>, Aohan Zeng<sup>1,4</sup>, Jianwei Zhang<sup>2</sup>, Runxin Zhong<sup>1</sup>, Tianhui Shi<sup>1</sup>, Sha Liu<sup>3</sup>, Weimin Zheng<sup>1</sup>, Jie Tang<sup>1,4</sup>, Hongxia Yang<sup>2</sup>, Xin Liu<sup>3</sup>, Jidong Zhai<sup>1</sup>, Wenguang Chen<sup>1</sup>

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The choice of the random blocks are diferent for each "iteration" : dynamic structures Each block is dense : depending of the hierarchy of the machine, we may use vectorial hardware at the lower level.

for Video Captioning

HPC challenges Kevin Lin; Linjie Li; Chung-Ching Lin; Faisal Ahmed, Zhe Gan, Zicheng Liu, Yumao Lu, Lijuan Wang Microsoft

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## 1 - Generator of non-Hermitian matrices, from given spectrum



### Example

![](_page_26_Figure_1.jpeg)

(a) Dominant Clustered Eigenvalues: acceptance = 94%, max error =  $3 \times 10^{-2}$ 

![](_page_26_Figure_3.jpeg)

(b) Clustered Eigenvalues: acceptance = 100%, max error =  $7 \times 10^{-5}$ 

![](_page_26_Figure_5.jpeg)

![](_page_26_Figure_6.jpeg)

![](_page_26_Figure_7.jpeg)

March

![](_page_27_Figure_0.jpeg)

![](_page_28_Picture_0.jpeg)

## 2 – Distributed and Parallel Generator of brain-scale graph-matrices

As we are speaking about "brain scale", we generate a graph as close as possible from the brain structure (to be updated with data from several researches – Neurospin/CEA)

Graphs, insipered form the topolgy of the humain brain

- Aprox. 10<sup>11</sup> Neurons,
- Up to 10 000 (different) conections,
- Several parts (left, rigth, entiric,...),
- Several feature per neuron.

It is also a possible to generate such sparse matrices for other kind of experiments

Graphs (sparse matrices) : several densities of nodes-neurons on each part, and several densities of connection from one part to another one (to be set in the future from data coming from RMI brain topology researches).

If we add some features to each neuron, we have a data set which may be adapted for **Graph Convolutional Network analysis, and others approaches** 

Size of the data set :  $10^{11} \times 10^4 + 10^{11} \times 10^{11}$  x number of features

It would be very expensive to upload the data (I/O), to experiment with several hypothesis : we have to generate the data directly in Parallel without I/O, using sparse adjaceny non-symetric matrices)

![](_page_29_Picture_12.jpeg)

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## Brain

## Parts of the brain

- part name
- % of connection to the opposite side
- % of connection to other parts
- number of neuron types in each part
- total number of neurons

### - Neurons

- neuron name
- number of neurons for each type of neuron
- number of connections for these neurons

#### + features

Very sparse and large non symetrical matrices

The connection inside each part, and betwwen parts are parmatrized and randomly set, for the moment, waiting more data.

Internal	Connections	Connections
connections	from part 2	from part 3
in part <b>1</b>	to part 1	to part 1
<u>Connections</u>	Internal	Connections
from part <b>1</b>	connections	from part 3
to part <b>2</b>	in part <b>2</b>	to part 2
Connections	Connections	Internal
from part <b>1</b>	from part 2	connections
to part <b>3</b>	to part 3	in part <b>3</b>

## + dense rectangular matrix for the features (n \* number features)

March 6, 2023

HPC challenges

PageRank? : ranking of the neurons (personalized PageRank would analyse the topological impacts of on set of neurons with the others)

Data set for Graph Convolutional Networks : to avoid oversmoothing, we drop edge and/or nodes. The ranking of the nodes (topologie – random walk based – **Pagerank**) may be use to optimize the dropping (method RankedDrop, IEEE Big Data 2022, Tokyo)

We first expriment on not too large graphs with such structures, the generation of the graphs and the ranking of the nodes, on different plateforms and supercomputers.

The % of connections between neurons-nodes are for the moment not based on the state-or-art, we are just evaluating our algorithms.

We don't compare with version with I/O as it is not relevant as it would be too expensive, Even if the matrices are not to large . We don't want to consume such energy.

On the future, we would like to run on exascale supercomputer to generate really brainscale graphs, and run (personalized) PageRank methods as examples.

How do neurons operate on sparse distributed representations? A mathematical theory of sparsity, neurons and active dendrites<sup>1</sup>

Subutai Ahmad<sup>1</sup>\*, Jeff Hawkins<sup>1</sup>

<sup>1</sup>Numenta, Inc., Redwood City, CA, USA

HPC challenges

## What is BTIDG2?

**BTIDG2**, short for Brain Topology Inspired Distributed Graph Generator, is able to generate very large graphs which is inspired by the topology of humain brain. This matrix generated is implemented based C, and parallised based on MPI.

As a software, **BTIDG2** ensures that:

- the generator must be able to work for large sparse matrices without any I/Os
- the generator must work in a distributed fashion
- support multiple (COO and CSR) distributed sparse data formats
- the generator must be as configurable as possible
- a **\*converter**\*, allowing to convert a data file containing the brain information as we found it on the internet into an input file for our matrix generator.

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BTIDG2 github repository: https://github.com/SMG2S/BTIDG2

View page source

BTIDG2 website: https://smg2s.github.io/BTIDG2/

#### License

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✤ » BTIDG2's Documentation

### **Brain Structure**

#### Structs

#### struct BrainPart

Structure containing the information of a part of the brain.

It contains :

a. The number of neuron types that can be encountered in the part b. the cumulative distribution of neurons in the neuron types

c. the probability of connection (for each type) to other parts of the brain.

This structure depends on the brain to which it belongs.

#### **Public Members**

#### int nbTypeNeuron

the number of neuron types that can be encountered in the part

#### double \*repartitionNeuronCumulee

the cumulative distribution of neurons in the neuron types

#### double \*probaConnection

the probability of connection (for each type) to other parts of the brain

#### struct Brain

Structure containing the information of a brain.

It contains :

- a. the total number of neurons
- b. the number of parts in the brain
- c. the indices (neurons) at which parts start
- d. the brain parts (see BrainPart).

#### **Public Members**

#### long long dimension

the total number of neurons

#### int nb\_part

the number of parts in the brain

#### long long \*parties\_cerveau

the indices (neurons) at which parts start

#### BrainPart \*brainPart

the brain parts (see BrainPart).

## void brainAdjMatrixCSR(csr \*M\_CSR, MatrixDistBlockInfo, Brain \*brain, int \*neuron\_types, BrainMatrixInfo \*debugInfo)

Generates a CSR square adjacency matrix of dimension (\*brain).dimension, corresponding to the brain passed as a parameter, in a two-dimensional process grid. A row/column of the matrix corresponding to a neuron, adjacency matrix means that it is the row-neuron that connect to the column-neuron

Condition : BlockInfo must have been filled with information suitable for the brain (Otherwise, the generation will not necessarily fail, but the generated matrix will not necessarily correspond to the brain)

#### Parameters

- [in] BlockInfo : structure containing information about the local mpi process ("block")
- [in] brain : Pointer to the brain, basis for the generation of the matrix
- [in] neuron\_types : vector containing the types chosen for the neurons of the brain passed as a parameter
- [out] M\_CSR : Pointer to a structure corresponding to a CSR matrix. At the end of the generation, contains the generated matrix.
- [out] debugInfo : OPTIONAL Pointer to a debug structure or NULL. If not NULL, at the end of the generation, contains debug information such as the number of connections made per neuron (<=> number of 1 on each row), the total number of connections, etc.

#### int get\_nb\_neuron\_brain\_part(Brain \*brain, int part)

Function that returns the number of neurons in a specific brain part.

Function that returns the number of neurons in the brain part if index part, in the Brain brain

Return

number of neurons in the brain part

Parameters

- [in] ind : index of the neuron
- [in] brain : pointer to a brain

#### void printf\_recap\_brain(Brain \*brain)

Brain print.

Function that displays a summary of the brain passed as a parameter (useful for debugging a brain)

#### Parameters

• [in] brain : pointer to a brain

3 - The convergence depend of the spectrum. we may analyze the efficiency of the different preconditoners with respect to the spectra, fro very large sparse matrices

![](_page_39_Figure_1.jpeg)

Figure: Convergence Comparison using a matrix generated by SMG2S.

HPC challenges

![](_page_40_Figure_0.jpeg)

![](_page_40_Figure_1.jpeg)

![](_page_40_Figure_2.jpeg)

(d) Spectral Distribution IV: matrix size = 2000,  $m_g = 20$ , d = 10. UCGLE(*eigen*<sub>1</sub>) has  $3 \times$  speedup, and UCGLE(*eigen*<sub>2</sub>) has  $2 \times$  speedup.

Xinzhe Wu , Serge G. Petiton, Yutong Lu: A parallel generator of non-Hermitian matrices computed from given spectra. Concurr. Comput. Pract. Exp. 32(20) (2020)

## Expriments on GRID5000 (France), JEREDA (Julich), and Fugaku (Kobe)

### Using only 256 cores , to generate not too large sparse matrices

![](_page_41_Figure_2.jpeg)

## Hardware Configuration of the JURECA DC Module (Phase 2: as of May 2021)

#### • 480 standard compute nodes

- 2× AMD EPYC 7742, 2× 64 cores, 2.25 GHz
- $\circ~$  512 (16× 32) GB DDR4, 3200 MHz
- InfiniBand HDR100 (NVIDIA Mellanox Connect-X6)
- diskless
- 96 large-memory compute nodes
  - 2× AMD EPYC 7742, 2× 64 cores, 2.25 GHz
  - 1024 (16× 64) GB DDR4, 3200 MHz
  - InfiniBand HDR100 (NVIDIA Mellanox Connect-X6)
  - diskless

#### 192 accelerated compute nodes

- 2× AMD EPYC 7742, 2× 64 cores, 2.25 GHz
- 512 (16× 32) GB DDR4, 3200 MHz
- $\,\circ\,$  4× NVIDIA A100 GPU, 4× 40 GB HBM2e
- 2× InfiniBand HDR (NVIDIA Mellanox Connect-X6)
- diskless

Block CSR NNZ = 3,5 % of N<sup>2</sup> Densities

• 0.5% inside each parts

**ATOS-BULL** 

5% between partsDiferent numbers of blocks

### We expriment with respect to several paramters :

- The matrix size
- The number of core
- Grid size and others pameters for the graph-matrices

#### JEREDA DC, Julich

#### **Process Grid tests**

We applied the following tests to matrices whose dimension vary from 30k to 80k, while keeping the same number of allocated resources (256 cores)

![](_page_42_Figure_3.jpeg)

![](_page_43_Figure_0.jpeg)

#### Minimum Generation time depending on the matrix dimension

Minimum Generation time depending on the matrix dimension

#### Matrix size tests

The following test were carried out by varying the size (dimension) of the matrix, while working with the same allocated resources (256 cores)

![](_page_44_Figure_2.jpeg)

Minimum Generation time depending on the matrix dimension

Minimum generation time depending on the matrix dimension

HPC challenges

#### **Strong scaling experiments**

In this experiment, we increase the number of allocated core while keeping the same problem size (a matrix with a dimension equal to 65536)

![](_page_45_Figure_2.jpeg)

Minimum Generation Time - Strong Scaling

Minimum generation time depending on number of allocated core

#### GRID5000

![](_page_46_Figure_2.jpeg)

Minimum total time (generation and pagerank) depending on the grid structure

#### Pagerank : seqence of sparse matrix-vector products + reduction\_with\_add

![](_page_47_Figure_0.jpeg)

Minimum total time (generation and pagerank) depending on the grid structure (zoom on "little matrices")

### Generation time :

![](_page_48_Figure_1.jpeg)

Minimum generation time depending on the grid structure

### PageRank time :

![](_page_49_Figure_1.jpeg)

Minimum PageRank time depending on the grid structure

![](_page_50_Figure_0.jpeg)

Minimum PageRank time depending on the grid structure (zoom on "little matrices")

First experiments done by Maxence Vandromme, part of a collaboration with Mitsuhisa Sato a and Miwako T

![](_page_51_Figure_1.jpeg)

#### Fugaku - PageRank - Brain generator v1

# compute nodes

GFlop/s

## Outline

- Introduction
- New levels of programming (Graphs of Tasks, Network on chip)
- New methods and algorithms (Unite&Conquer, Stochastic Matrix,..)
- HPC and Machine Learning (GCN, Transformer,..)
- Generators of Data Sets and matrices for brain-scale applicatons
- What post-exascale plateforms and programing paradigms

## Conclusion

We propose two generators of data to be able to expriment "brain-scale" applications without expensive I/O, both for CSE and Machine Learning ; directly compute in parallel.

These generators may also be used to preconditione or pre-train (transfert learning) methods, while we upload other data, if we have enough knownledge of the spectrum or of the topology of the targeted graphs-matrices

### HPC challenge and new computing frontiers

- Arithmetic : mixed, new normalisations?
- New methods : optimisation of operations, of iterations-epochs (unite&conquer or Neural-Network ensembles), minimization of communications,
- Hierarchical architecture : cluster-cloud-distributed + parallelism + NOC chips, (accelerated) set of cores,
- Progamming paradigms : graph of task, PGAS-data parallelism, vectorial
- (Non-Hermitian) Sparse linear algebra : sequence of matrix-vector products, sequence of (dynamic) sparse matrix products, eigenvalues
- New applications : "brain scale" bigbird transformer, AI, human brain, ...